
Priority Inheritance Protocols: An Approach to Real-Time Synchronization

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Priority Inversion Problem

❑ In traditional synchronization techniques

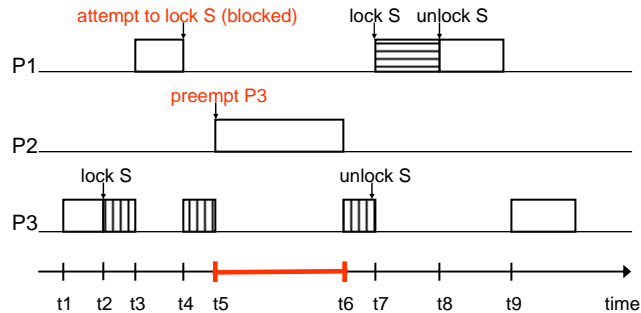
- Higher priority tasks can be blocked by lower priority tasks

❑ Motivating example

- Three tasks: P1, P2, P3
 - P1 has the highest priority
- P3 locks the semaphore S
- P1 preempts P3 and then attempts to lock S, but is blocked
 - This necessary for mutual exclusion
- P2 preempts P3, delaying P1

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Priority Inversion Example

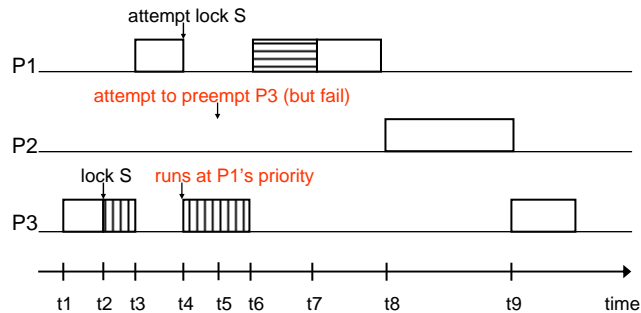


- ❑ The blocking of P1 is affected by independent P2
 - In the interval (t5, t6)
- ❑ The blocking duration is unpredictable

Priority Inheritance Protocol

- ❑ A simple, but not good solution
 - A task in its critical section is not allowed to be preempted
 - But, higher priority tasks may suffer unnecessary blocking
- ❑ Priority Inheritance Protocol
 - When a task blocks one or more higher priority tasks
 - It ignores its original priority
 - It executes its critical section at the highest priority level of all the tasks it is blocking
- ❑ Priority inheritance is transitive
 - If P3 blocks P2, and P2 blocks P1, P3 inherits P1's priority

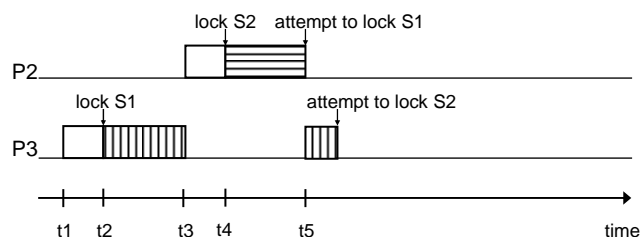
Priority Inheritance Example



- ❑ P3 inherits the P1's priority from t4
- ❑ P2 cannot preempt P1 at t5

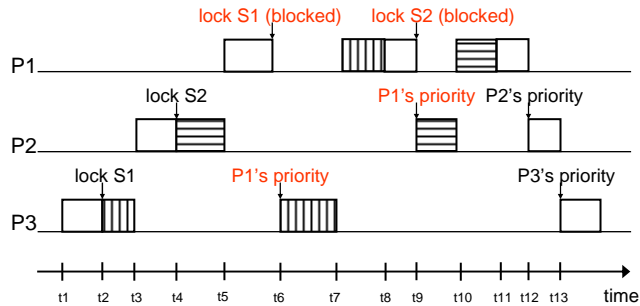
Limitations of PIP

- ❑ PIP have two problems
 - Do not prevent deadlocks
 - Chain of blocking can be formed



**P2 and P3 are in nested critical sections
 resulting in deadlocks!**

Chained Blocking Example



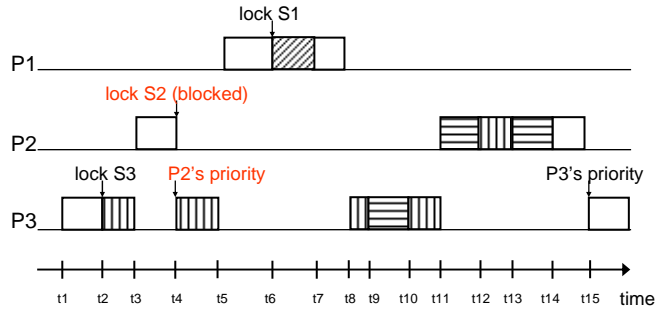
❑ P1 is blocked two times

- When it attempts to lock S1
- When it attempts to lock S2

Priority Ceiling Protocol

- ❑ Based on Priority Inheritance Protocol
- ❑ A priority ceiling is assigned to each semaphore
 - The highest priority of tasks that may use the semaphore
- ❑ Task P is allowed to start a new critical section
 - Only if P's priority is higher than all priority ceilings of all the semaphores locked by jobs other than P
- ❑ PCP can
 - Prevent deadlocks
 - Avoid chained blocking

PCP Example



- P1: {S1}, P2: {S2, S3}, P3: {S3, S2}
- At t3, P2's priority is higher than P3
- At t4, P2's priority is not higher than S3's ceiling
- At 6, P1's priority is higher than P2's priority and S3's ceiling